Mobile Systems M



Alma Mater Studiorum – University of Bologna CdS Laurea Magistrale (MSc) in Computer Science Engineering

> Mobile Systems M course (8 ECTS) II Term – Academic Year 2022/2023

01 – Overview on Wireless Communications

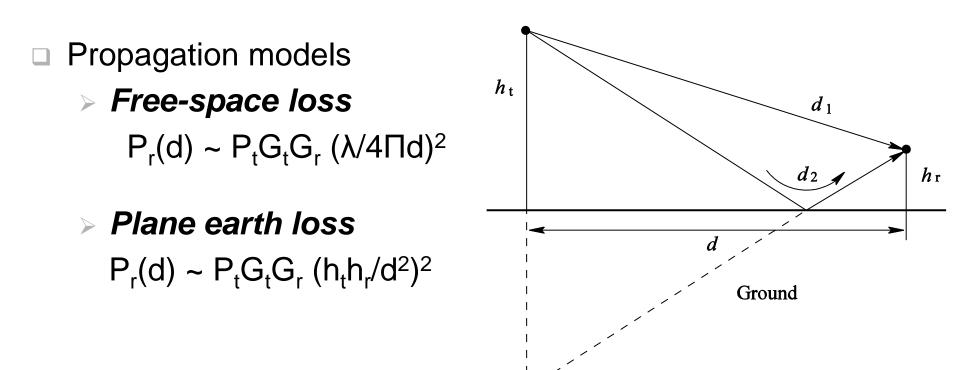
Paolo Bellavista paolo.bellavista@unibo.it



A few elements of Wireless Propagation Models

Exactly to make you worried immediately from the beginning ③, all and all always comes from Maxwell's equations...

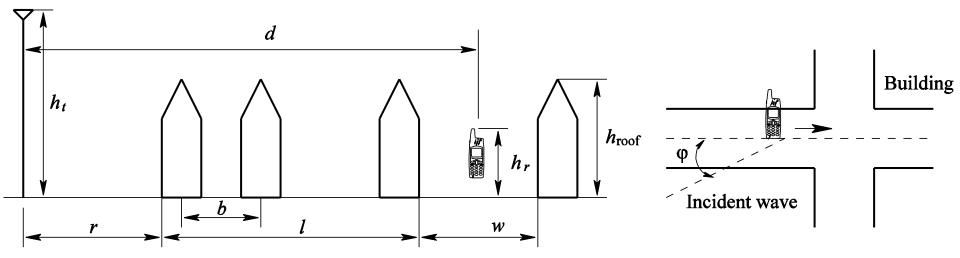
You did discuss a lot about propagation and fading models in the course «Fondamenti di Telecomunicazioni T», didn't you?





A few elements of Wireless Propagation Models

- In *real or realistic* scenarios, actual conditions are far more complex; this pushes, as very frequently occurs in engineering, towards *approximated empirical models*
 - Propagation models
 - Plane earth loss
 - $\mathsf{P}_{\mathsf{r}}(\mathsf{d}) \sim \mathsf{P}_{\mathsf{t}}\mathsf{G}_{\mathsf{t}}\mathsf{G}_{\mathsf{r}} \; (\mathsf{h}_{\mathsf{t}}\mathsf{h}_{\mathsf{r}}/\mathsf{d}^2)^2$

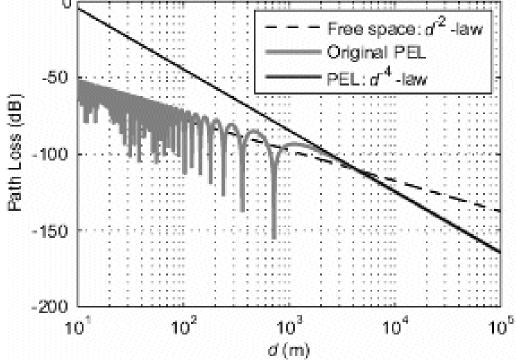




A few elements of Wireless Propagation Models

In real or realistic scenarios, actual conditions are far more complex; this pushes, as very frequently occurs in engineering, towards *approximated empirical models*

- Propagation models
 - > **Plane earth loss** $P_r(d) \sim P_tG_tG_r (h_th_r/d^2)^2$
 - > Okumura-Hata
 - > COST-231-Hata
 - COST-231-Walfisch-Ikegami
 - > ...
 - Indoor propagation models
 - Ray-tracing-based propagation models (geometric optics and diffraction)



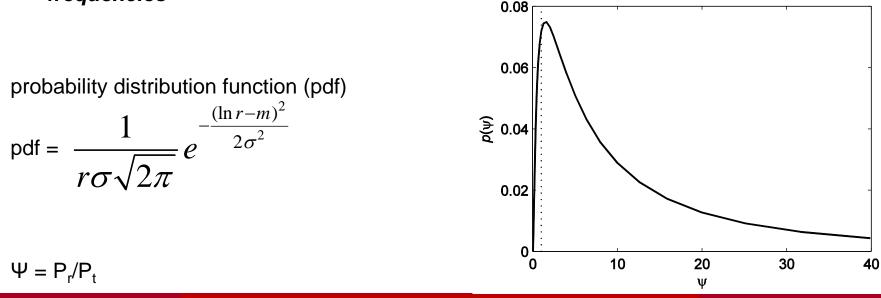


A few elements of Channel Fading Models

Three primary types of channel fading in mobile communications (physically due to reflection, diffraction, and scattering):

- 1. Shadowing (or slow-fading)
- 2. Multipath *Rayleigh* fading
- 3. Frequency-selective fading

Shadowing: total absorption and partial reflection by trees, buildings, mobile vehicles, ... decrease of received power strength in a wide spectrum of frequencies





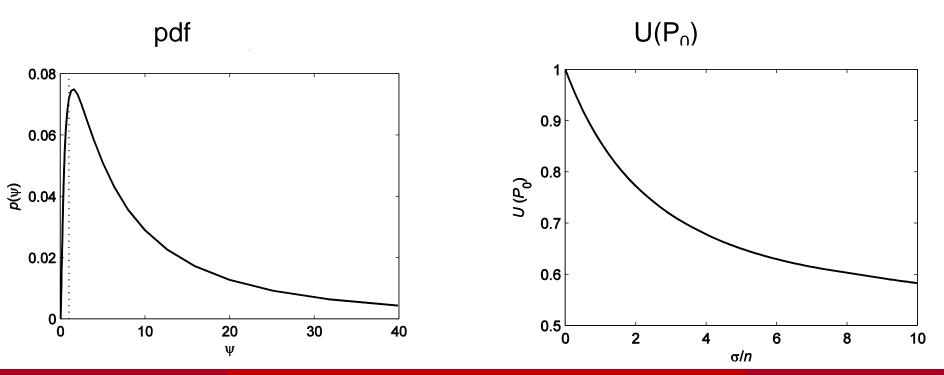
Shadowing: Probability!

Shadowing:

from *pdf* to the determination of *U(P₀)*, that is

probability of being in the coverage area

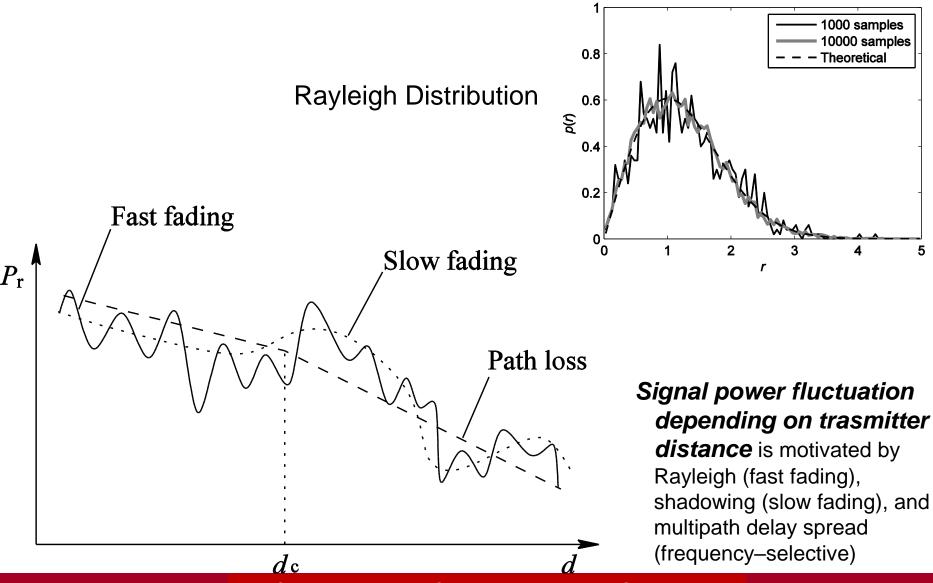
(percentage of the coverage area with received signal power strength greater than P_0)



Overview on Wireless Communications – Mobile Systems M

Rayleigh Fading and Combined Effects

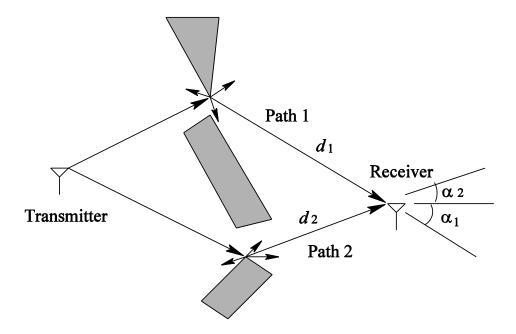




Rayleigh Fading and Combined Effects



It is possible to demonstrate that Rayleigh fading can be explained as the result, for example, of a simple *multipath model, with 2 paths*



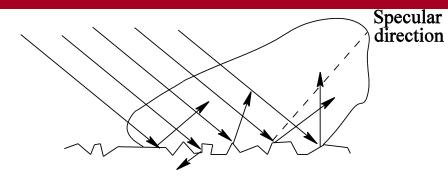
In addition, **Doppler effect** and frequency shift (time-selective fading)

But also Ricean fading, Nagakami fading, Suzuki fading, ...

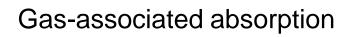
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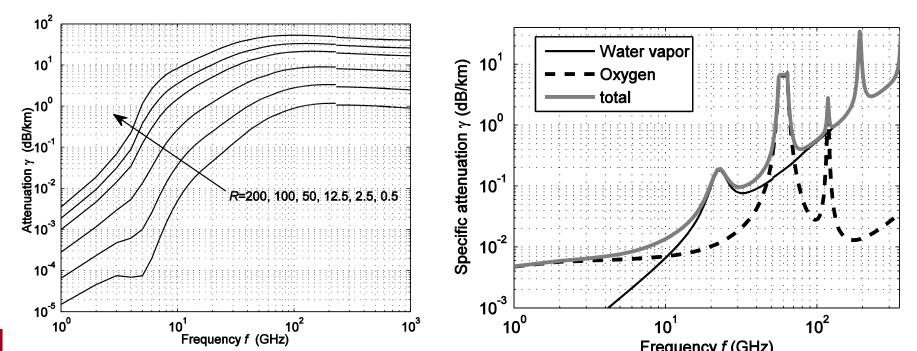
Additional Effects: Scattering, Athmosperic Absorption, ...

Scattering due to rough surface



Rain-associated attenuation







- THE THE PARTY OF T
 - Complexity of real world
 - Probabilistic approaches
 - Only partial math/physical models
 - Empirical models
 - Need for in-the-field measurements

Often, no need to be very sophisticated in these aspects because of the *many and many simplifying assumptions* imposed by any model...

Pragmatic approach

Anyway, do not worry *©*, this course will not discuss anymore about the physical layer



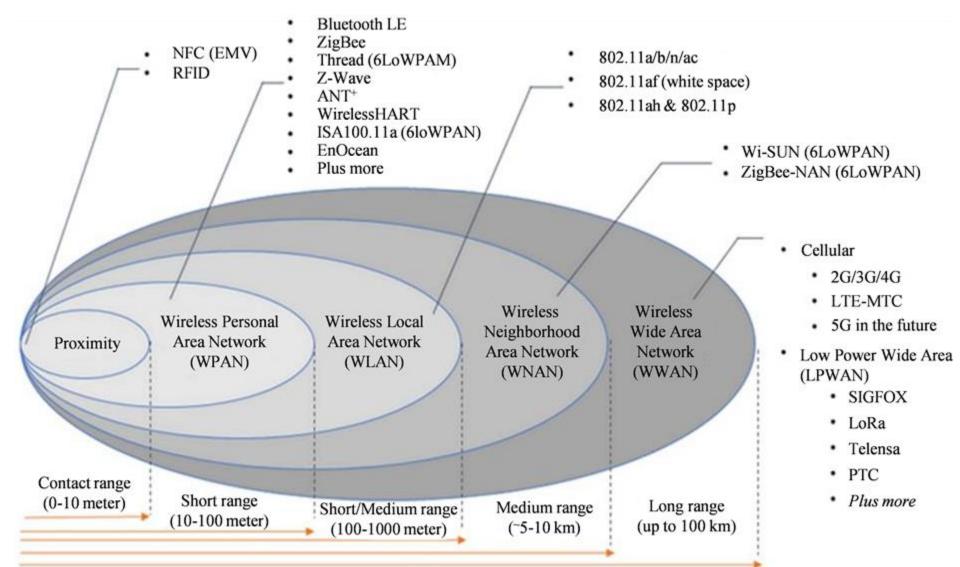
Wireless Communications: Which Connectivity Technologies today?

Overview on:

- □ Wireless LAN IEEE 802.11
- Wireless MAN
 - > IEEE 802.16 WiMAX/WirelessMAN
 - > IEEE 802.20/802.11p MBWA/Vehicular Mobility
 - > IEEE 802.11 WiFi & Mesh Networking
- Cellular Networks
- □ Wireless PAN (IEEE 802.15)
 - > Bluetooth and Bluetooth Low Energy (BLE)
 - > ZigBee, Low rate WPAN, ...

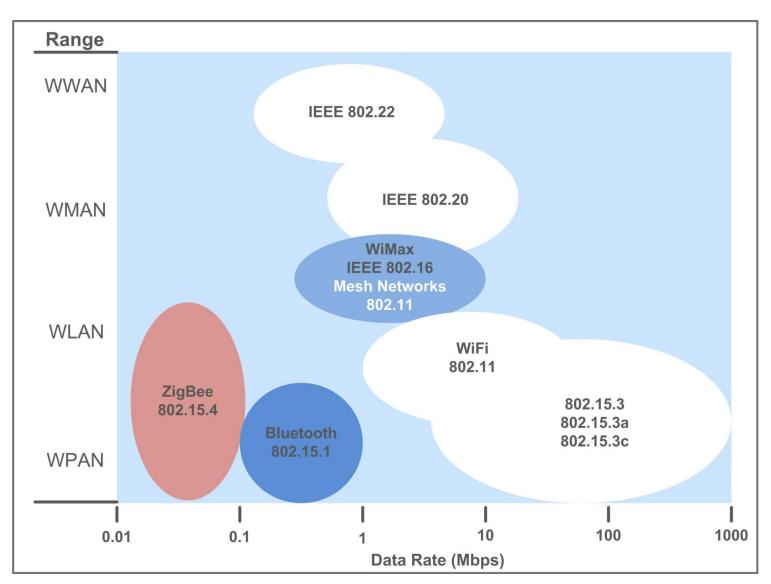


Wireless Space





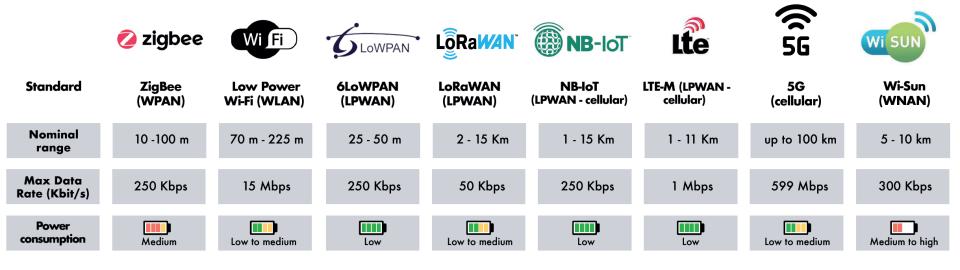
Wireless Communications Space: with finer Details...



Overview on Wireless Communications – Mobile Systems M



Wireless Communications Space: specialized for IoT...

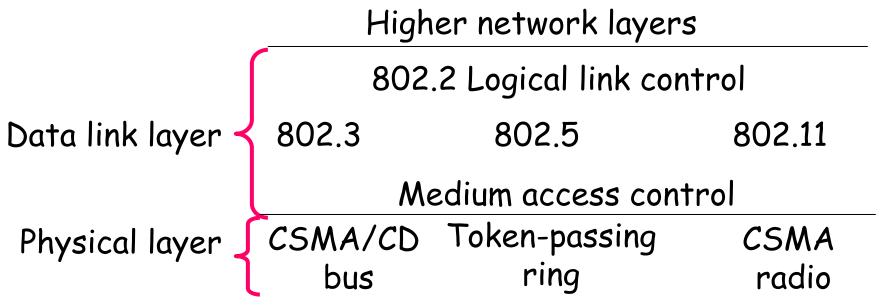




IEEE 802.x Standard for Local Area Networks (LAN)

The IEEE 802.x standards define:

- > Physical layer protocols
- Datalink layer protocols
 - Medium Access Control (MAC) sub-layer
 - Logical Link Control (LLC) sub-layer





TRADITIONAL CSMA/CD Algorithm:

> To perform channel sensing

- If idle, immediate transmission
- > If occupied, wait until the channel becomes idle again

Collision Detection

- Immediate abort of a transmission if collision detected
- Following transmission try after waiting for a random time interval

Why CSMA/CD?

- > CSMA relevantly reduces the number of collisions
- > CD reduces the negative impact of potential collisions

OPTIMISTIC APPROACH



Exactly the same Ethernet Algorithm

1-persistent CSMA + binary exponential backoff

Channel sensing

- □ If *idle*, immediate transmission (*p*-persistent, here p=1)
- Otherwise, wait until the channel is idle and then immediate transmission
- If collision is detected
 - Random selection of a transmission slot in the range [0, 2ⁿ), where n=min(10,k) and k is the number of collisions already occurred for that frame
 - Slot length = 2 * end-to-end propagation delay
 - > Transmission during the selected slot
- Algorithm self-adapts well to the dynamic changes in network load





- Wireless communications for so-called short range
 - > Transmission range < 250m</p>
- Bandwidth up to 866Mbps (with many and many variants)
 - > **802.11b** (many variants, up to 2 and 11Mbps at 2.4GHz)
 - > 802.11a (many variants, up to 54Mbps at 5GHz)
 - > 802.11g (many variants, up to 54Mbps at 2.4GHz)
 - 802.11n (many variants, up to 150Mbps at 2.4GHz)
 - 802.11ac (many variants, up to 866Mbps at 5GHz)
 - 802.11ax (WiFi-6, 2019, many variants, up to 1201Mbps)
 - > **802.11be** (under discussion, 2024, WiFi-7, 40Gbps?)
- Standardization of MAC layer, of physical layer, and of some security aspects (partial support)
 - > Typically no Quality of Service



From several perspectives, the IEEE 802.11 approach is similar to Ethernet, but:

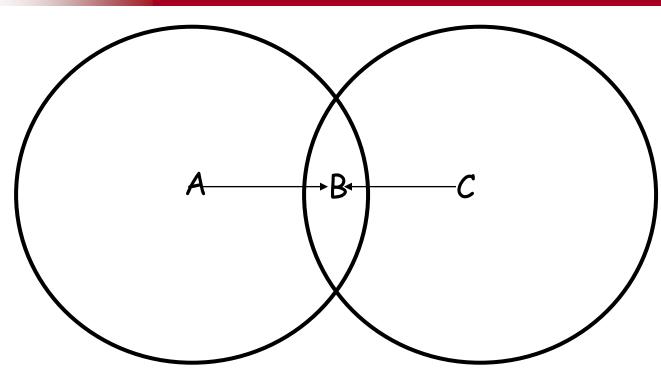
- □ Can CSMA/CD work correctly over wireless medium?
- Simple forms of CD are NOT possible (may fail) because of attenuation issues
- Also CS exhibits some issues: not all nodes are in the coverage range of any other node

□ What affected by *mobility*?

- Let's start with "classical" NEW problems, which are not possible in wired communications:
 - > Hidden node
 - > Exposed node

The so-called Hidden Node Issue

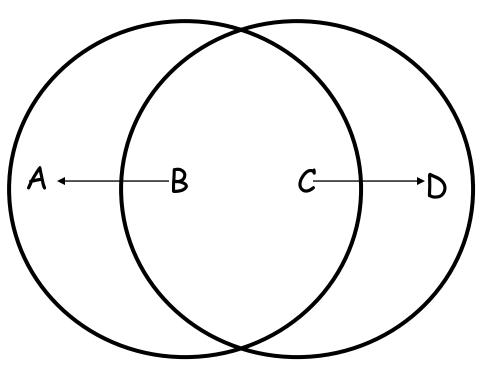




- □ A starts transmitting to B
- C cannot be aware of A's communication (out of range) and therefore starts transmitting to B
 Collision!

The so-called **Exposed Node Issue**

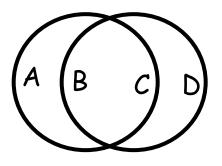




- B starts transmitting to A
- C detects this transmission and does NOT send any communication towards D
- Error: this communication would NOT have interfered with A and D



Multiple Access Collision Avoidance (MACA)



B is willing to transmit to C

- B transmits a first *Ready To Send* (RTS) frame towards C. The frame includes the length of data to be transmitted
- C replies with a Clear To Send (CTS) frame
- Other nodes overhear the RTS frame and HAVE to stay inactive (in silence) for sufficient time to allow C to reply with its CTS
- Other nodes overhear the CTS frame and HAVE to remain in silence for the whole duration of B's transmission



- Collisions may occur ANYWAY
 - For instance, simultaneous sending of RTS frames
 - But at least data frames CANNOT collide
 - > And because of *mobility?*

Solution: nodes anyway make re-transmissions by using the "classical" Ethernet algorithm, i.e., binary exponential backoff

Try to countercheck experimentally how often the RTS/CTS mechanism is really employed in your WiFi networks?How? For instance, via sniffing of the radio channel...Ever used *WireShark* or similar?



IEEE 802.11: Two Primary Configurations

- Ad hoc: all nodes are potentially mobile and communicate directly the ones with the others. We will see in MANETs...
 - How many hops? Which availability and reliability? Which latency is still acceptable? Which trade-offs between those elements?
 - Which issues/challenges if implemented with the MAC/routing/transport mechanisms and algorithms that you already know?

Differences wrt Wi-Fi Direct?

Wi-Fi Direct devices host sw-based access points ("Soft APs")...

Base station:

- Fixed nodes, equipped also with wired connectivity, serve as access points (or base stations) for wireless communications
- Mobile nodes communicate through access points
- Mobile nodes operate occasional scans in their locality to find new/other access points by using a simple probe/response prot
- Which issues/challenges if implemented with the MAC/routing/transport mechanisms and algorithms that you already know? For instance, dynamic transfer from one access point to one another



Wi-Fi Direct, or Wi-Fi P2P (initial name),

enabler for *direct connection* of off-the-shelf devices, also *multi-point*, competitor of Bluetooth for applications that do not need low energy consumption

Only one of the Wi-Fi devices HAS to be compliant with Wi-Fi Direct to establish a peer-to-peer connection

Description Pairing

We'll return to it after Bluetooth and MANET presentations

Wi-Fi Direct essentially FORCES any compliant device to implement a *software access point (Soft AP)*. Soft AP implements a version of the AP connection mode called *Wi-Fi Protected Setup* (*push-button or PIN-based setup*)

When a device enters the range of a Wi-Fi Direct host, it obtains setup info through Protected Setup in simplified and quasi-automated way



THE REPORT

A Soft AP may be implemented according to different levels of complexity, depending on the role to be played (standard *flexibility*)

For instance, a digital picture frame can provide only very base services for connectivity and upload by cameras; a *smart phone typically offers more complex data tethering* with bridge capability towards the Internet

Note that (virtualization vs physical location):

"Pairing" of Wi-Fi Direct devices may be configured to request *real proximity of devices*, for example, via verification of Near Field Communication (NFC) or Bluetooth signals (and/or explicit button press by the devices' users)

(Brief Note on) WiFi Direct Soft AP



Are WiFi Direct nodes peer-to-peer (no infrastructured) nodes in a cooperative network?

More info in the following parts of the course...



Overview on Wireless Communications – Mobile Systems M





- Goal: broadband access in metropolitan areas
- Why?
 - Growing demand for *local network connectivity with high bandwidth*
 - High-capacity fiber/cable towards any user is typically unsustainable from the economic point of view (think to WiMAX in mountain areas or for covering geographic areas with limited population density)
 - > IEEE 802.11 is short range (around 250m in optimal conditions)
 - Cellular communications (see next lectures) could be or could be perceived as too expensive for the offered datarate



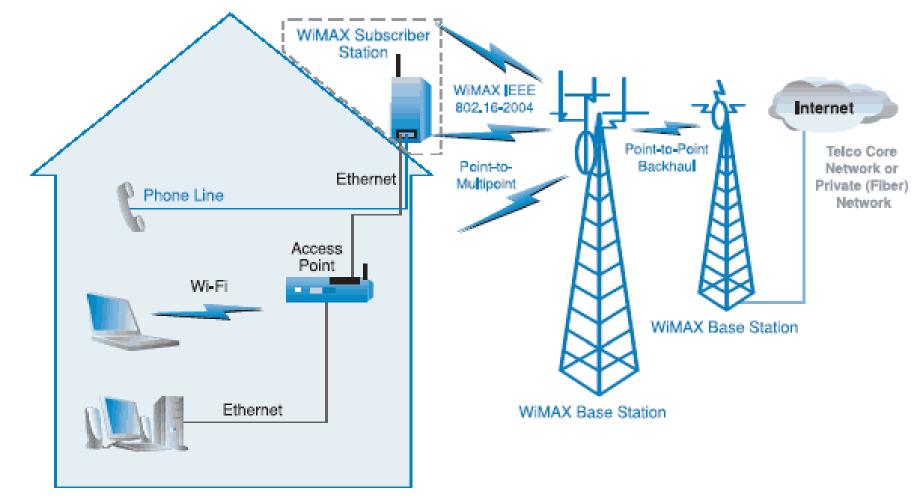
Wireless MAN: IEEE 802.16

□ IEEE 802.16 standard (WiMAX) *to cover "last mile"*

- Maximum datarate = 70Mbps (in first version; think to current datarates of IEEE 802.11 variants and 5G cellular networking)
- Optimized for "fixed stations"
 - Basic idea: every building equipped with an antenna (subscriber station) that communicates with central base stations, wired connected to the Internet
 - > Circuit commutation, hierarchical architecture
- In your opinion, why *limited spread*, in particular in some national contexts (e.g., Italy)?

Wireless MAN: IEEE 802.16





Customer Premise (Home, Business or HOTSPOT)



Wireless MAN: WiFi Coverage?

- School of Engineering and Architecture, Viale Risorgimento, relatively well covered [©] by WiFi APs
- □ How to properly cover more extensive regions?
 - > Already done in several cities, of various sizes

Brooklin Massachusetts, Rochelle Illinois, Thorold Ontario, Milton Keyens UK, Oulu Finland, Galatsi Greece, and many others are examples of municipalities with the need to provide WiFi access for city internal operations, public safety, meter reading and more, ...

City-wide WiFi mesh networks, IEEE 802.11s, for example interested students can see C. Gomez, J. Paradells, "Urban Automation Networks: Current and Emerging Solutions for Sensed Data Collection and Actuation in Smart Cities", MDPI Sensors 2015, doi:10.3390/s150922874

Open question: how to organize and manage a WiFi mesh network? Which technological challenges?



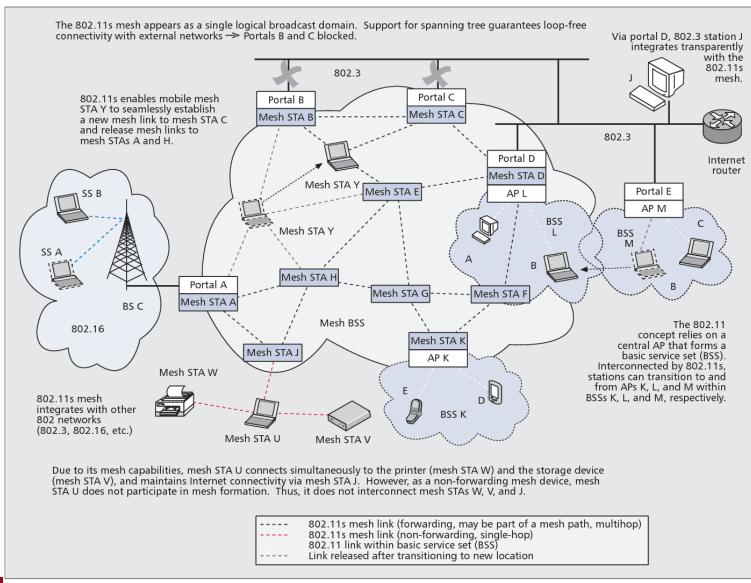
IEEE 802.11 for Municipal Mesh Network

IEEE 802.11 base stations compose an infrastructure and can wirelessly communicate the one with the other

- Compose a mesh (mesh networking)
- Can manage client connection handoffs between different base stations (*support to mobility*)
- > Which technical differences wrt "regular" IEEE 802.11?
- Why mesh-based solution?
 - Costs, robustness, …
 - Possibility of connecting only a few base stations to the Internet via wired connections
- Ideally any domestic or university campus AP (or even any mobile node) can *contribute to the mesh infrastructure*
 - > Which problems? Which needs?

IEEE 802.11 for Municipal Mesh Network







Mobile Broadband Wireless Access (MBWA): IEEE 802.20 – Vehicular Mobility

IEEE 802.20 - ABANDONED - standardized physical and MAC layers for *mobile broadband wireless access (MBWA) systems*

- Worked in licensed band under 3.5 GHz
- > Optimized for IP-oriented data communications
- Datarate peek in downlink per user > 1 Mbps
- Datarate peek in uplink per user > 300 Kbps
- Number of concurrent active users HIGHER than in other mobile wireless solutions
- Supported different classes of vehicular mobility, up to 250 Km/h in MAN environments
- But which other issues are present in terms of routing, transport, application-layer, ...?

NOW IEEE 802.11p? ...



Basically IEEE 802.11 variant *with no need of long association time to AP* (basic service set – BSS)

IEEE 802.11p is an approved amendment for *wireless access in vehicular environments (WAVE)*

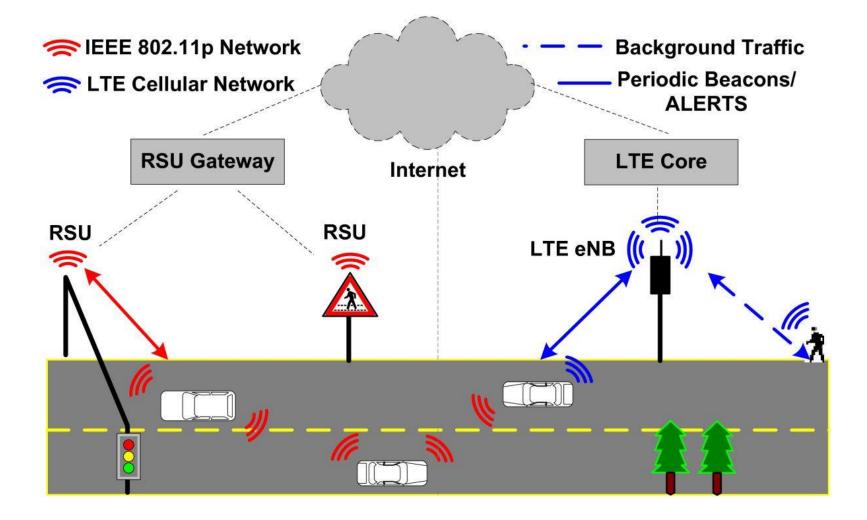
Includes data exchange between *high-speed vehicles* and between the *vehicles and the roadside infrastructure*

Licensed ITS band of 5.9 GHz (5.85-5.925 GHz)

IEEE 1609 is a higher layer standard based on 802.11p and also the base of a European standard for vehicular communication known as *ETSI ITS-G5*



IEEE 802.11p – Vehicular Mobility

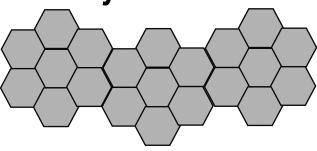




Cellular Networking: Basic Idea

Many transmission antennas with low power (800-900MHz)

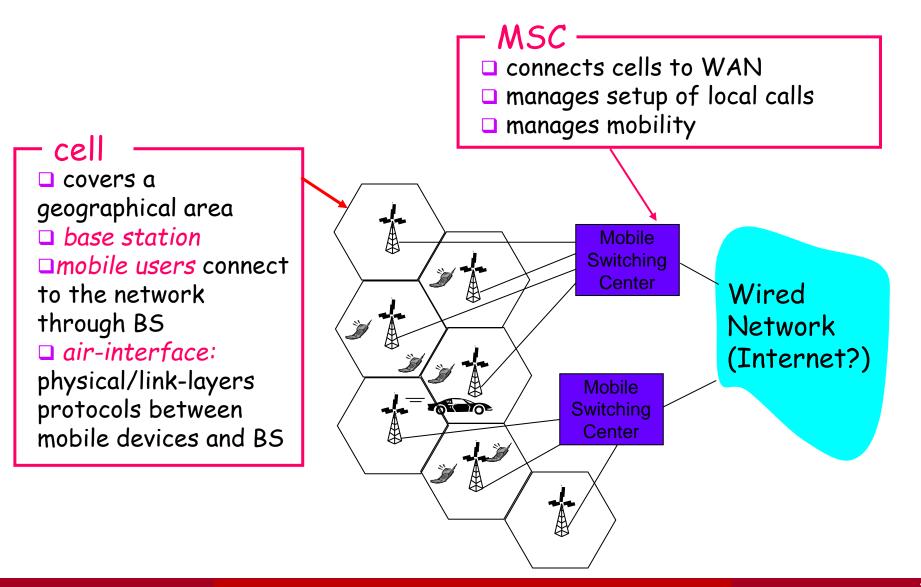
- > Why? Scalability and incremental deployment
- □ Geographical area split into *adjacent cells*



- Any cell is served by one base radio station (Base Station - BS)
 - Analogously to what seen for IEEE 802.11 BS (access point)
- Radio frequencies can be re-used in non-adjacent cells, provided that a proper spatial interval (of non-interference) is respected



Cellular Network Architecture





A Plethora of Different Standards for Cellular Networking (1)

2G systems: only voice channels (but already digital tech in both communication and commutation)

- > IS-136 TDMA: FDMA/TDMA combined (N.America)
- Global System for Mobile communications (GSM): FDMA/TDMA combined
 - Most widespread and used
- > IS-95 Code Division Multiple Access (CDMA)

2.5G systems: channels for voice and data

- Extensions to 2G before the widespread adoption of 3G
- General Packet Radio Service (GPRS)
 - □ GSM evolution
 - Data sent over multiple channels, when available



Cellular Networking Standards (2)

□ ... 2.5G

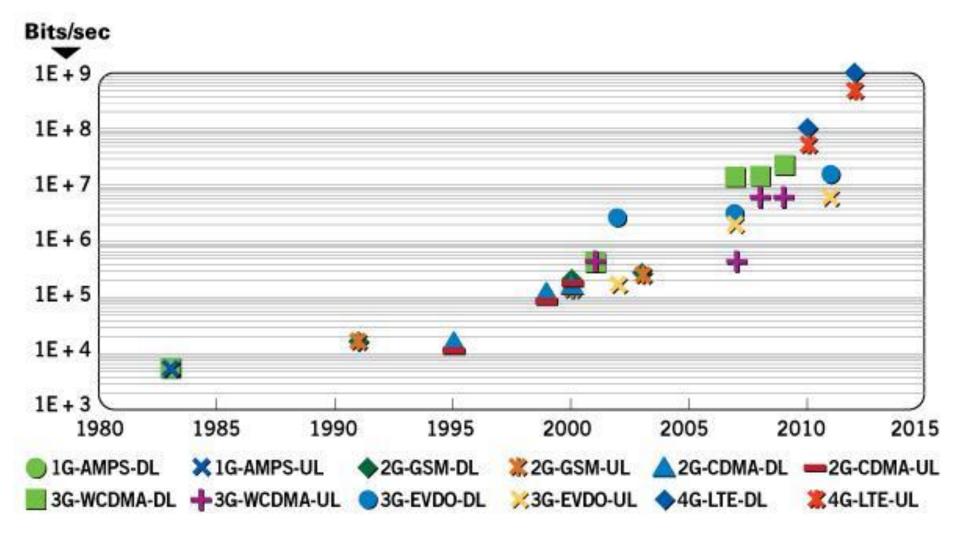
- Enhanced Data rates for Global Evolution (EDGE)
 - Also in this case, GSM evolution; modulation change
 - Datarate up to 384Kbps
- 3G systems: voice and data channels, of course
 - > Universal Mobile Telecommunications Service (UMTS)
 - □ Further evolution step of GSM, but with CDMA exploitation
 - CDMA-2000

□ **4G systems**: voice and data channels, of course

- Datarate upto 1000Mbps (download) and 500Mbps (upload)
- > Often indicated as LTE or LTE-Advanced
- Already available in Japan since 2006

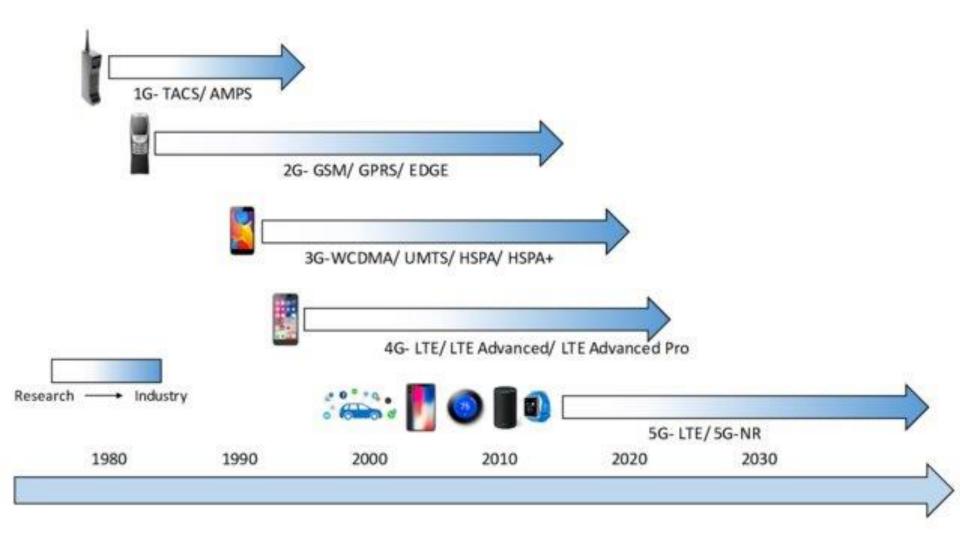


Cellular Networking Standards: Evolution

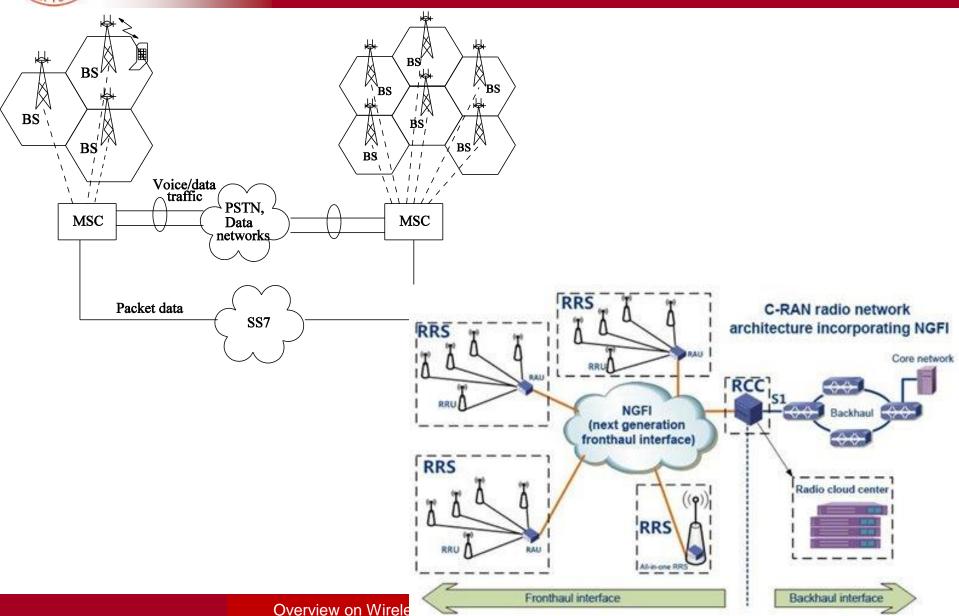




Cellular Networking Standards: Evolution

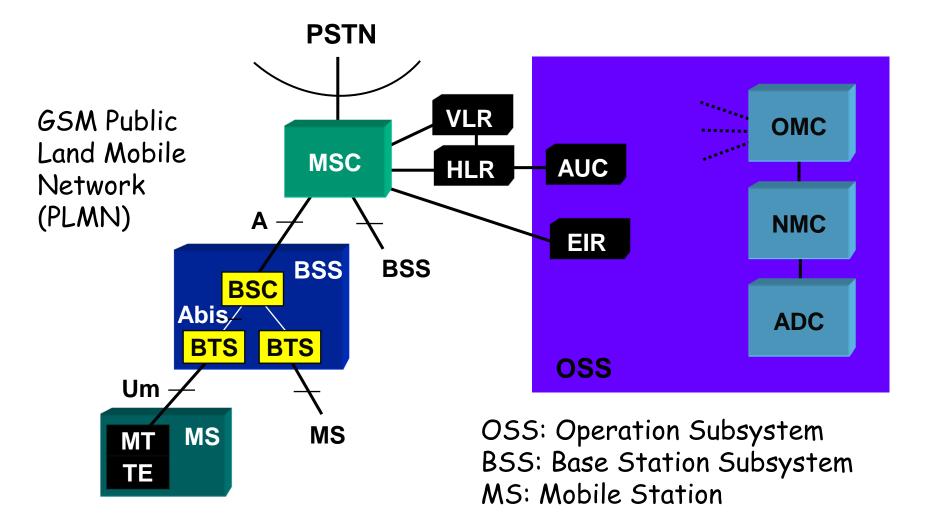


Cellular Networking: from Circuit to Packet Switching



GSM: General Architecture (1)







- MS communicates with several BSSs via its radio interface Um
- Mobile Terminal (MT) supports physical channels
 between MS and BSS (radio transmission, channel coding, voice conversation coding, ...)
- Terminal Equipment (TE) contains terminal/userspecific data (also user profiles) and smart card data (e.g., associated with SIM)
 - Identifies the user towards the network, e.g., when supporting personal mobility (in addition to terminal mobility) and security

By the way, do you know the differences among *user mobility*, *terminal mobility*, *session/resource mobility*?



Hierarchical organization and locality principle

- BSS communicates with MSC via network interface A
- Base Transceiver Station (BTS) manages channel allocation, signaling, frequency hopping, handover triggers, …
- BTS communicates with Base Station Controller (BSC) by using the Abis interface
- BSC manages radio channels, paging, handoff for different BTS
- □ MSC plays the role of *gateway to PSTN and to packet data networks*
 - Works for switching, paging, location updates of MS, handoff control, …
- Home Location Register (HLR) stores subscribers' info and part of the location data of MS to perform ingress call routing towards the proper Visitor Location Register (VLR)
- VLR stores the list of "visiting" users in its area and assigns the socalled roaming numbers

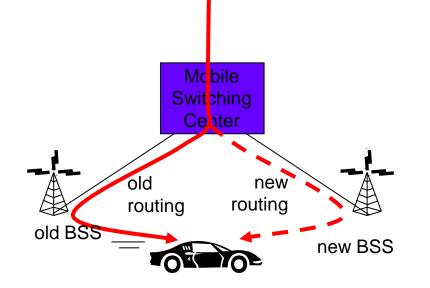


- AUthentication Center (AUC) is accessed by HLR for users' authentication before service access
- Equipment Identity Register (EIR) supports the identification of stolen or fraudolent (stolen identity) MS
- OSS works for network monitoring, control, and management
 - Operations and Maintenance Center (OMC)
 - Network Management Center (NMC)
 - > ADministration Center (ADC)

GSM Handoff (under the same MSC)

- Motivations for handoff (or handover)
 - Stronger signal from/to new BSS ("more continuous" connectivity, minor battery consumption, ...)
 - Load balancing: it may free channels for the current BSS
 - GSM does NOT specify why to operate handoff (policy), but only how (mechanisms)

Handoff triggered by the OLD BSS







- 1. Old BSS informs MSC of the need to perform handoff, by providing it with a *list of one or more new BSSs*
- 2. MSC *configures the new path (resource allocation)* towards the new BSS
- 3. New BSS *allocates the radio channel* that the mobile "visitor" will have to use
- 4. New BSS signals to MSC and old BSS signals its *ready state*
- 5. Old BSS informs mobile device of *the need to operate handoff towards the new BSS*
- 6. It is the mobile device that signals to new BSS to activate the new channel
- 7. Mobile device signals to MSC, via the new BSS, that the *handoff has been completed; MSC performs call re-routing*
- 8. *Resources are released* on the MSC-to-OldBSS path

new BSS

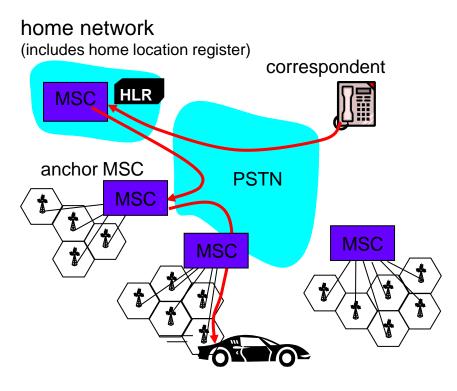
Mobile

Switching Center

old BSS





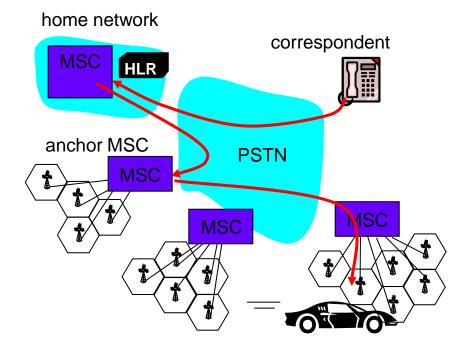


(a) before the handoff

- Anchor MSC: first visited MSC during the call
 - Call remains routed through the anchor MSC
- New MSCs are appended at the end of the MSC chain, step by step, as the mobile device is moving
- IS-41 allows *minimizing the path* (*optional feature*) in the *multi-MSC chain*





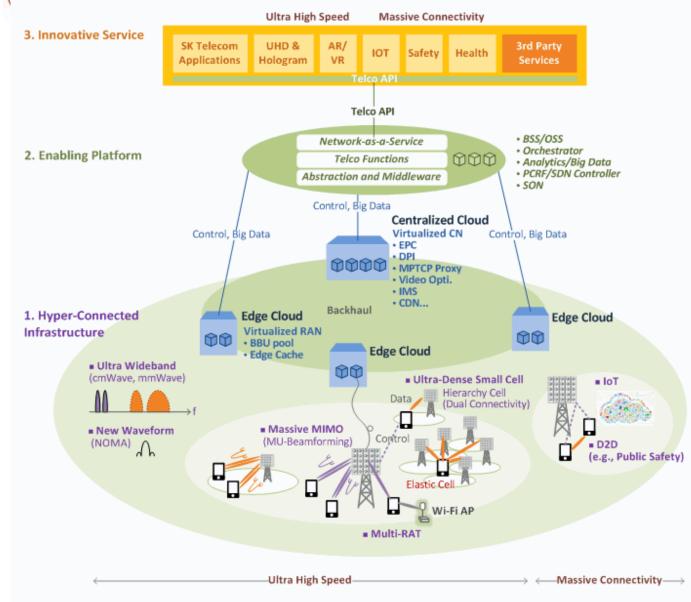


⁽b) after handoff

- Anchor MSC: first visited MSC during the call
 - Call remains routed through the anchor MSC
- New MSCs are appended at the end of the MSC chain, step by step, as the mobile device is moving
- IS-41 allows *minimizing the path* (*optional feature*) in the *multi-MSC chain*



Need for Clarifications about



As you may see, ALSO:

5G?

- Virtualization
- Middleware
- Cloud
- Edge cloud

We'll discuss 5G and mobile edge technologies in the last lectures...



Several taxonomies and classifications are possible (see also the wide literature in the field)

- Horizontal (homogeneous connectivity) or vertical (het connectivity) handoff
- Mobile-initiated or network-initiated. In addition, where is the handoff decision taken?
- □ *When* to trigger the handoff procedure?
 - > Reactive
 - > Proactive
- □ How to choose the *destination network*?
 - Which *monitoring indicators* to considerate and with which *metrics*?

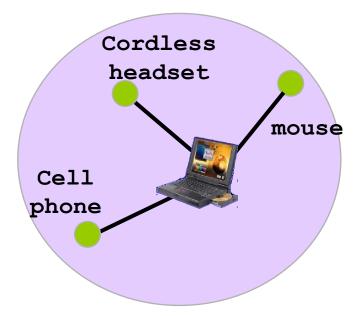


- Hard handoff or soft handoff management (to concurrently enable multiple «virtual» radio interfaces or only one?)
 - > Which pros and cons?
- Even multiple paths, also heterogeneous...
 - > Which technical challenges? What is currently supported?
- Standardization efforts and *IEEE 802.21* Media Independent Handover Services (for both 802 and non-802 networking solutions)



Wireless PAN: Bluetooth or IEEE 802.15.1

- Radio interface (2.4 GHz band) used for the connection and communications of wireless devices
 - Designed as technology for cable replacement
 - It could (?) be used for multi-hop communications in ad hoc networks
 - Short-range: 10-100 m
 - "Traditional" peek datarate: 1 Mbps
 - Low cost: goal was \$5-\$10 per interface at the proposal (widely achieved and passed nowadays)
 - Manages both voice and data







Bluetooth was born in 1998 from the alliance of IBM, Ericsson, Intel, Nokia, and Toshiba (*Bluetooth SIG*)

Objective: to design and implement a short-range radio communication system, based on small-size and low-cost components

- Bluetooth SIG defined the specifications of a set of *protocols* for radio communications and the associated *software stack*
- In addition to basic specifications, also definition of profiles that describe how to use the protocols by guaranteeing intercommunication between het devices

Bluetooth



- 2.4 GHz communications (same band of IEEE 802.11) in frequency hopping
- Initial theoretical bandwidth of 1306 kb/s (v2.0 + EDR)
- Consumed power: 1mW–10 mW
- Piconet: network topology with one master and 7 slaves at maximum
- Scatternet

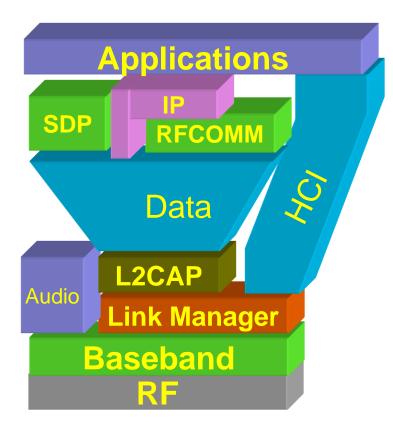


Differences wrt infrastructured IEEE 802.11:

- Point-to-point communications (between master and slaves)
- Needs discovery of devices and services
- Very limited broadcast mechanism



Protocols Stack (articulated...)



- Radio, Baseband, and Link
 Manager are usually
 implemented in firmware
 - The higher layers usually in software
 - Which tradeoff between implementations at different layers?
- Applications interact with firmware via *Host Controller Interface* (*HCI*)
 - Defined and supported interfaces: RS232, USB, PC Card

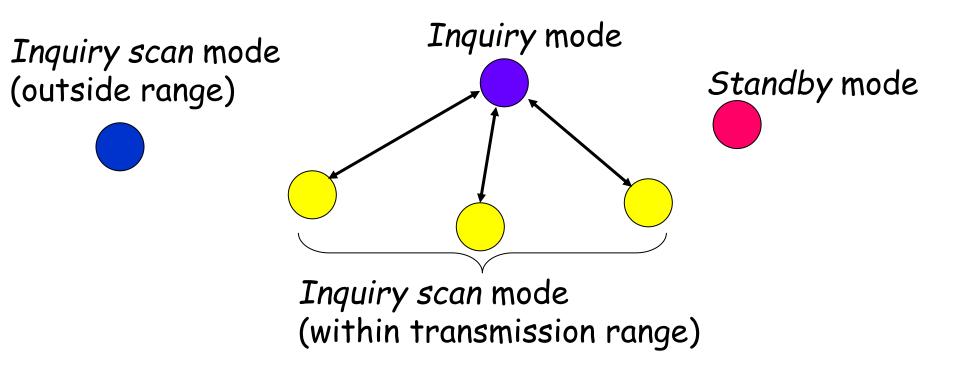


- Bluetooth specifies a protocol for the construction of the current topology implemented at the baseband layer
 - 1. **Inquiry phase**: a node in need of communication performs the **discovery of nearby nodes** (within its transmission range)
 - Who starts the protocol becomes the master, discovered nodes play the role of slaves
 - 2. Page/paging phase: master establishes a bidirectional communication channel (frequency hopping) with its slaves
 - Piconet is the base topology unit for Bluetooth
 - One master
 - Up to 7 active slaves (up to 256 slaves in parked state)
 - Single channel in frequency hopping



Inquiry Phase

Goal: to collect sufficient info on neighbor nodes in order to establish a piconet

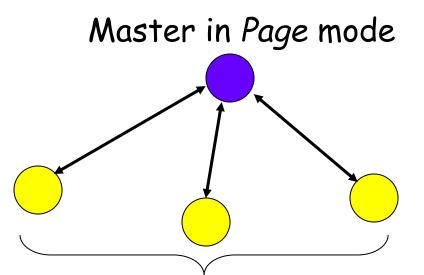




- Nodes that need to communicate enter the Inquiry mode state
 - They transmit their short ID (Inquiry Access Code IAC)
 - They wait for replies from neighbors
- To save energy, other devices can alternate Inquiry Scan mode and Standby mode
 - Inquiry Scan = listening to communication requests
 - Nodes reply when necessary
- After a given time interval, master enters the Page phase provided that it has received at least one reply

Page Phase





Slave in *Page scan* mode (within transmission range) Master knows the value of the estimated clock and BT_ADDR per each slave

Master performs broadcast
 of packets with ID

When a slave replies,
 master and slave go on with the exchange of the
 info needed to establish
 connection

Piconet communications are typically connectionoriented

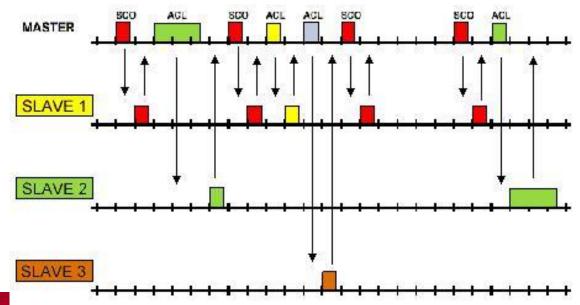


Bluetooth exploits frequency hopping in 2.4 GHz ISM

band (79 hop frequencies at 1 MHz distance). **Hopping sequence is** determined based on the master address; all devices in the piconet HAVE to follow the same sequence

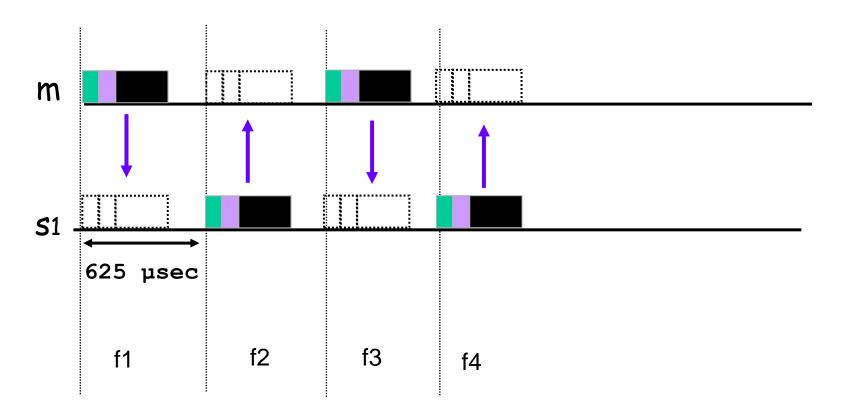
 Master communicates with each slave (rounds) in a periodical way Time is split into 625us slots

Master transmits at the beginning of odd slots, slaves at the beginning of even slots. *Need for synch at the piconet level*





Piconet Communications



Channel is split into time slots – each one of 625µs

Time Division Duplex: master and its slaves alternate in transmitting/listening



Bluetooth synchronizes most of its operations with *clock signal in real time*. For instance, to synch data exchange between devices, to determine which packets are re-transmitted or lost, to generate a pseudo-random sequence that is predictable and reproducible Bluetooth clock is implemented via a counter (28 bit) that is set to 0 when switching on the device and that increments every 312.5 μ s (half slot)

- > Any Bluetooth device has its own native clock (CLKN)
- CLK: piconet clock, coincides with the CLKN of the piconet master. All piconet nodes HAVE to synch their CLKN with CLK
- CLKE: also this clock derives from CLKN through offset; it is used by the master in the case of creation of a connection to a new slave, before that the slave is synch-ed with the master

A *master transmission always start when CLK[1:0] = 00 (odd indexed slots)*, while a slave transmission starts always when CLK[1:0] = 10 (even indexed slot)



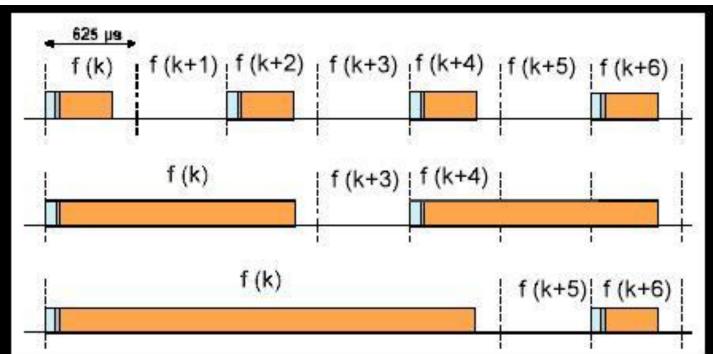
Bluetooth connections can be of two types: *connectionoriented and connectionless*, designed for the support of voice applications and data transfer: asynch service with no connection *(ACL, Asynchronous ConnectionLess)* and synch service oriented to connections *(SCO, Synchronous Connection Oriented)*

- ACL, data traffic and best-effort service ACL supports packet-switched, point-multipoint, and symmetric/asymmetric connections. Symmetric connections with maximum datarate of 433.9 kbit/s, asymmetric connections 723.2/57.6 kbit/s; a slave can transmit only if in the previous slot it received a packet by the master
- SCO, connections with real-time and multimedia traffic Circuit-switched, point-point, symmetric connections, voice transportation in 64 kbit/s channels, usually; master can support up to 3 SCO connections towards the same slave or different slaves in the same piconet

no re-transmission in the case of errors or losses (why?)

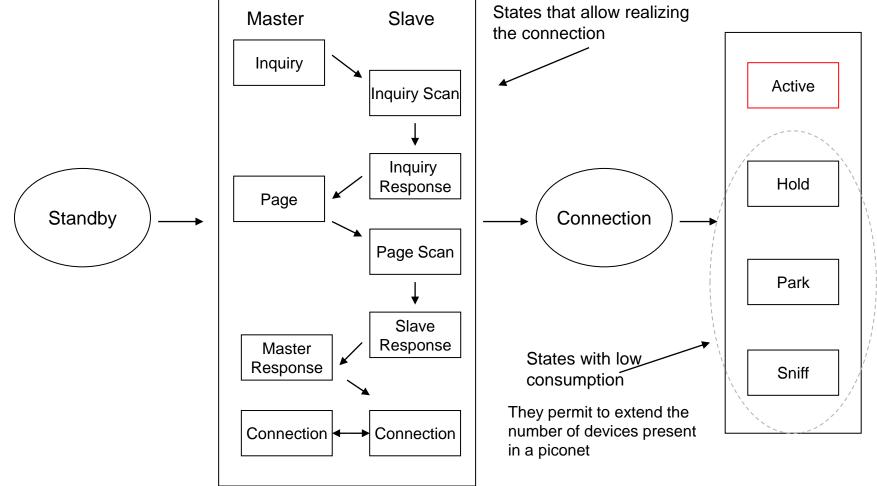


- In SCO connections, channel "reserved" for 2 time slots for communications between the master and one specific slave *Slot reservation periodicity is decided by the master (similar to circuit switching)*; time slots are allocated independently by the need of transmission
- ACL transmissions may occur in *pause intervals between SCO reserved slots*

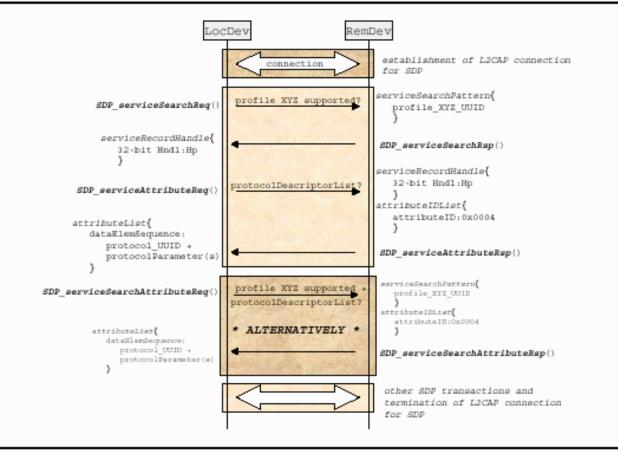


Device States





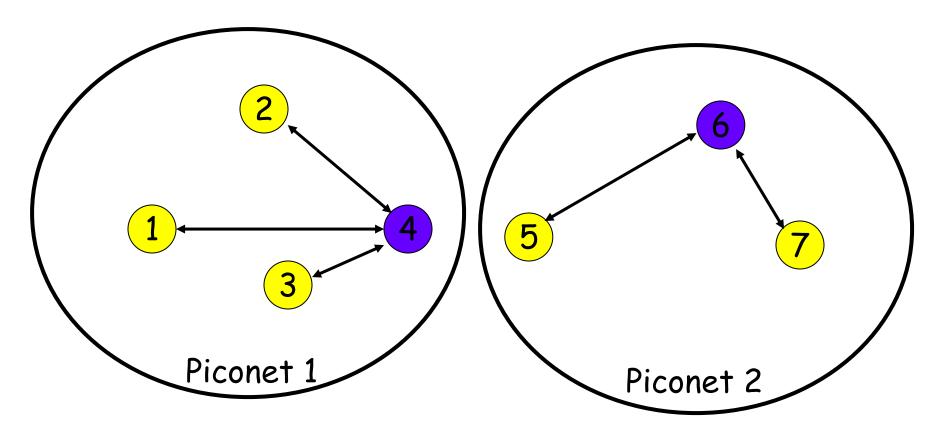
Also need to go up in the abstraction layer: Service Discovery Protocol (SDP)



- To discover available devices and services, Bluetooth applications exploit Service Discovery Protocol
- Any server keeps updated its database, available for clients, including the service records related to the offered services

Does Bluetooth Support Multi-Hop Communications?



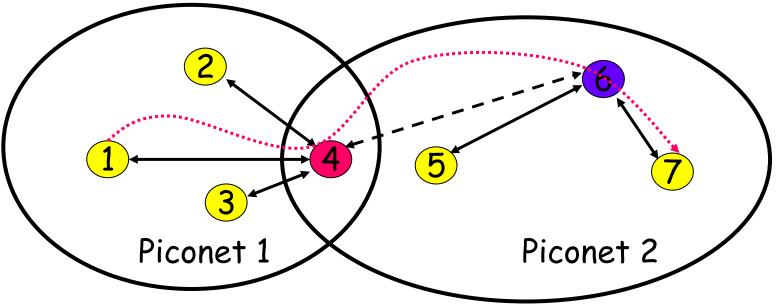


For instance, how can node1 communicate with node7 if the node distance is greater than their transmission range?



Bluetooth Scatternet

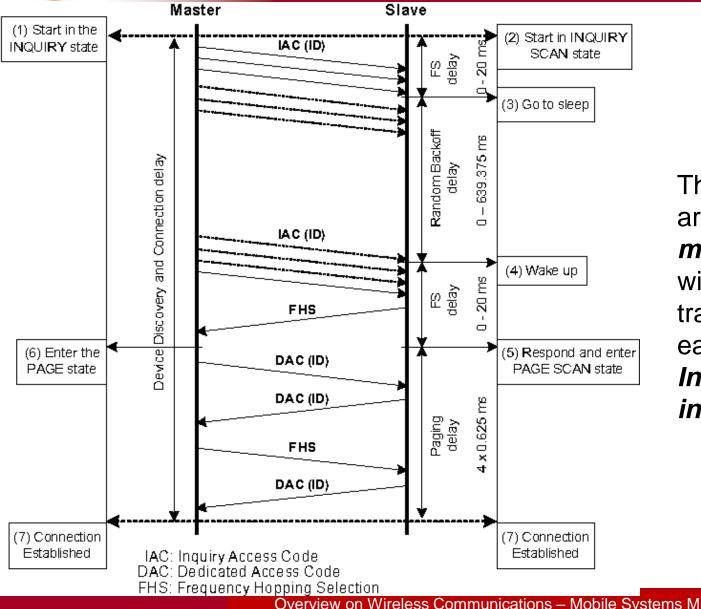
- Connection of multiple piconets, in proximity, via the sharing of either a common master OR of a common slave
- A node may be slave in a piconet and master in another one
 - > That node can "move" back and forth between the two piconets
- For a long time, *no commercial implementation* available (only a few academic experiments)



Node4 plays the role of master in Piconet1 and of slave in Piconet2
 A path may be created from node1 to node7 via nodes 4 and 6



Clashing Issues in Inquiry/Paging



The maximum values are measured when *multiple devices*,

within the transmission range of each other, *enter in Inquiry/Paging mode in the same interval*



Just to complete the overview...

- Compared to Classic Bluetooth, Bluetooth Smart (or BLE) is intended to provide considerably reduced power consumption and cost while maintaining a similar communication range
- Originally introduced under the name Wibree by Nokia in 2006; merged into the main Bluetooth standard in 2010 with the adoption of the Bluetooth Core Specification Version 4.0
- The Bluetooth SIG predicts that by 2022 more than 85% Bluetoothenabled devices will be BLE compliant
- BLE operates in the same spectrum range as Classic Bluetooth, but with a different set of channels (40 2-MHz channels). Within a channel, data is transmitted using Gaussian frequency shift modulation. Bitrate is 1 Mbit/s, and the maximum transmit power is 10 mW



BLE uses *frequency hopping* to counteract narrowband interference problems. Classic Bluetooth also uses frequency hopping but the details are different

Advertising and discovery

BLE devices are detected through a process based on *broadcast advertising packets*. This is done using 3 separate channels (frequencies), in order to avoid interference Both the advertising devices, and the scanners looking for those devices, step through the three channels in sequence only about a 1 in 9 chance of a scanner detecting a particular device's advertisement time to discover devices highly variable and long

Software model

All Bluetooth Smart devices (or BLE devices) use the Generic Attribute Profile (GATT)

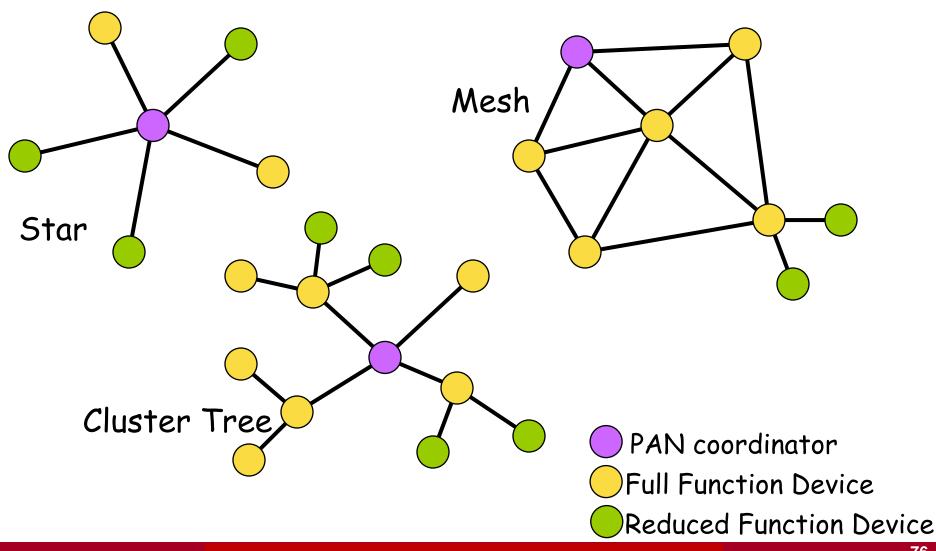


Wireless PAN: ZigBee or IEEE 802.15.4

- This standard is specifically designed for sensors and actuators networks with reliability, cost-effectiveness, and low power requirements
- Applications:
 - Control of home appliances
 - > Automation and control of large buildings and industrial plants
 - Monitoring of patients and elders, in particular at homes
 - Environmental monitoring
 - War applications and homeland security
- As Bluetooth, ZigBee provides *profiles* for higher layers of the support/application stack
- □ Up to 65536 nodes (clients) in a network
- Optimized for time-critical apps and for limited power consumption
 - Join < 30ms; state change from sleep to active < 15ms</p>
- Support to full mesh networking

ZigBee Topologies







Types of Roles and Channel Access Options

- Different types of roles for ZigBee devices
 - Coordinator (one and only one for each ZigBee network)
 - Starts the network formation
 - Can serve as *router* once the network is operating
 - "Full" device
 - Participates to message routing
 - > Device with limited features/functions
 - Only sensing and actuating operations, no routing
- Options for channel access
 - Non-beaconed network
 - Exploits CSMA-CA, almost the same way as already described
 - Positive acks for successful reception of packets
 - Beacon-enabled network
 - Coordinator transmits beacons at regular time intervals
 - Dedicated bandwidth and low latency
 - Low power mode also for the coordinator