Introduction Number Theory

Background

We will use a bit of number theory to construct:

- Key exchange protocols
- Digital signatures
- Public-key encryption

Notation

From here on:

- N denotes a positive integer.
- p denote a prime.

Notation: $\mathbb{Z}_N = \{0, 1, \ldots, N-1\}$

Can do addition and multiplication modulo N

Modular arithmetic

Examples: let $N = 12$

$9 + 8 = 5$ in $\mathbb{Z}_{12}$

$5 \times 7 = \square$ in $\mathbb{Z}_{12}$

$5 - 7 = \square$ in $\mathbb{Z}_{12}$

Arithmetic in $\mathbb{Z}_N$ works as you expect, e.g $x \cdot (y+z) = x \cdot y + x \cdot z$ in $\mathbb{Z}_N$

Modular arithmetic

Examples: let $N = 12$

$$9 + 8 = 5 \quad \text{in } \mathbb{Z}_{12}$$

$$5 \times 7 = 11 \quad \text{in } \mathbb{Z}_{12}$$

$$5 - 7 = 10 \quad \text{in } \mathbb{Z}_{12}$$

Arithmetic in $\mathbb{Z}_N$ works as you expect, e.g $x \cdot (y+z) = x \cdot y + x \cdot z$ in $\mathbb{Z}_N$

Greatest common divisor

<u>**Def**</u>: For ints. x,y: **gcd(x, y)** is the <u>greatest common divisor</u> of x,y

Example: gcd(12, 18) = 6

<u>**Fact**</u>: for all ints. x,y there exist ints. a,b such that

$$a \cdot x + b \cdot y = gcd(x,y)$$

a,b can be found efficiently using the extended Euclid alg.

If gcd(x,y)=1 we say that x and y are <u>**relatively prime**</u>

<u>**Example:**</u> 2 x 12 -1 x 18 = 6

Modular inversion

Over the rationals, inverse w.r.t. the moltiplication of 2 is ½ .
What about $\mathbb{Z}_N$?

<u>Def</u>: The **inverse** of x in $\mathbb{Z}_N$ is an element y in $\mathbb{Z}_N$ s.t. $x \cdot y = 1$

y is denoted x^{-1} .

Example: let N be an odd integer.

The inverse of 2 in $\mathbb{Z}_N$ is $\frac{N+1}{2}$ since $2 \cdot \frac{N+1}{2} = N + 1 = 1$

Modular inversion

Which elements have an inverse in $\mathbb{Z}_N$?

Lemma: x in $\mathbb{Z}_N$ has an inverse if and only if gcd(x,N) = 1

Proof:

gcd(x,N)=1 $\Rightarrow$ $\exists$ a,b: $a \cdot x + b \cdot N = 1 \Rightarrow a \cdot x = 1$ in $\mathbb{Z}_N$

$\Rightarrow x^{-1} = a$ in $\mathbb{Z}_N$

gcd(x,N) > 1 $\Rightarrow$ $\forall$a: gcd($a \cdot x$, N) > 1 $\Rightarrow$ $a \cdot x \neq 1$ in $\mathbb{Z}_N$

More notation

<u>Def:</u> $\mathbb{Z}_N^*$ = (set of invertible elements in $\mathbb{Z}_N$) =

= { x∈ $\mathbb{Z}_N$: gcd(x,N) = 1 }

Examples:

1. for prime p, $\mathbb{Z}_p^* = \mathbb{Z}_p \setminus \{0\} = \{1, 2, \ldots, p-1\}$
2. $\mathbb{Z}_{12}^*$ =

For x in $\mathbb{Z}_N^*$, can find x^{-1} using extended Euclid algorithm.

More notation

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Examples:

1. for prime p, $\mathbb{Z}_p^* = \mathbb{Z}_p \setminus \{0\} = \{1, 2, \ldots, p-1\}$
2. $\mathbb{Z}_{12}^*$ = { 1, 5, 7, 11}

For x in $\mathbb{Z}_N^*$, can find x^{-1} using extended Euclid algorithm.

Solving modular linear equations

Solve: **$a \cdot x + b = 0$ in** $\mathbb{Z}_N$

Solution: **$x = -b \cdot a^{-1}$ in** $\mathbb{Z}_N$

Find a^{-1} in $\mathbb{Z}_N$ using extended Euclid. Run time: $O(\log^2 N)$

What about modular quadratic equations?

next segments

Fermat's theorem (1640)

<u>Thm</u>: Let p be a prime

$$\forall\, x \in (Z_p)^* : \quad x^{p-1} = 1 \text{ in } Z_p$$

Example: p=5. $3^4 = 81 = 1$ in Z_5

Example of application:

So: $x \in (Z_p)^* \quad \Rightarrow \quad x \cdot x^{p-2} = 1 \quad \Rightarrow \quad x^{-1} = x^{p-2}$ in Z_p

another way to compute inverses, but less efficient than Euclid

Application: generating random primes

Suppose we want to generate a large random prime

say, prime p of length 1024 bits (i.e. $p \approx 2^{1024}$)

Step 1: choose a random integer $p \in [\ 2^{1024}\ ,\ 2^{1025}-1\ ]$

Step 2: test if $2^{p-1} = 1$ in Z_p

If so, output p and stop. If not, goto step 1 .

Simple algorithm (not the best). **$\Pr[\ p \text{ not prime }] < 2^{-60}$**

The structure of $(Z_p)^*$

<u>Thm</u> (Euler): $(Z_p)^*$ is a **cyclic group**, that is

$\exists\ g \in (Z_p)^*$ such that $\{1, g, g^2, g^3, \ldots, g^{p-2}\} = (Z_p)^*$

g is called a **<u>generator</u>** of $(Z_p)^*$

Example: p=7. $\{1, 3, 3^2, 3^3, 3^4, 3^5\} = \{1, 3, 2, 6, 4, 5\} = (Z_7)^*$

Not every elem. is a generator: $\{1, 2, 2^2, 2^3, 2^4, 2^5\} = \{1, 2, 4\}$

Order

For $g \in (Z_p)^*$ the set $\{1, g, g^2, g^3, \ldots\}$ is called

the **group generated by g**, denoted $<g>$

<u>**Def**</u>: the **order** of $g \in (Z_p)^*$ is the size of $<g>$

$$\mathbf{ord_p(g) \quad = \quad |<g>| \quad = \quad (smallest\ a>0\ s.t.\ \ g^a = 1\ in\ Z_p)}$$

Examples: $ord_7(3) = 6$; $ord_7(2) = 3$; $ord_7(1) = 1$

<u>**Thm**</u> (Lagrange): $\forall g \in (Z_p)^*$: **$ord_p(g)$** divides $p-1$

Euler's generalization of Fermat (1736)

Def: For an integer N define $\varphi(N) = |(Z_N)^*|$ (Euler's φ func.)

Examples: $\varphi(12) = |\{1,5,7,11\}| = 4$; $\varphi(p) = p-1$

For $N=p\cdot q$: $\varphi(N) = N-p-q+1 = (p-1)(q-1)$

Thm (Euler): $\forall x \in (Z_N)^*$: $x^{\varphi(N)} = 1$ in Z_N

Example: $5^{\varphi(12)} = 5^4 = 625 = 1$ in Z_{12}

Generalization of Fermat. Basis of the RSA cryptosystem

Modular e'th roots

We know how to solve modular **linear** equations:

$\mathbf{a \cdot x + b = 0}$ in Z_N Solution: $\mathbf{x = -b \cdot a^{-1}}$ in Z_N

What about higher degree polynomials?

Example: let p be a prime and $c \in Z_p$. Can we solve:

$$x^2 - c = 0 \quad , \quad y^3 - c = 0 \quad , \quad z^{37} - c = 0 \quad \text{in} \quad Z_p$$

Modular e'th roots

Let p be a prime and $c \in Z_p$.

<u>**Def**</u>: $x \in Z_p$ s.t. $x^e = c$ in Z_p is called an <u>**e'th root**</u> of c .

$6^3 = 216 = 7$ in $\mathbb{Z}_{11}$

Examples: $7^{1/3} = 6$ in $\mathbb{Z}_{11}$

$3^{1/2} = 5$ in $\mathbb{Z}_{11}$

$1^{1/3} = 1$ in $\mathbb{Z}_{11}$

$2^{1/2}$ does not exist in $\mathbb{Z}_{11}$

The easy case

When does **$c^{1/e}$ in Z_p** exist? Can we compute it efficiently?

<u>The easy case</u>: suppose gcd(e , p-1) = 1

Then for all c in $(Z_p)^*$: **$c^{1/e}$** exists in Z_p and is easy to find.

The case e=2: square roots

If p is an odd prime then gcd(2, p-1) $\neq$ 1

$$\begin{tikzcd}[column sep=small] x \arrow[dr] & & -x \arrow[dl] \\ & x^2 & \end{tikzcd}$$

<u>Fact</u>: in $\mathbb{Z}_p^*$, $x \longrightarrow x^2$ is a 2-to-1 function

<u>Example</u>: in $\mathbb{Z}_{11}^*$:

$$\begin{tikzcd}[column sep=small] 1 \arrow[dr] & & 10 \arrow[dl] & 2 \arrow[dr] & & 9 \arrow[dl] & 3 \arrow[dr] & & 8 \arrow[dl] & 4 \arrow[dr] & & 7 \arrow[dl] & 5 \arrow[dr] & & 6 \arrow[dl] \\ & 1 & & & 4 & & & 9 & & & 5 & & & 3 & \end{tikzcd}$$

<u>Def</u>: x in $\mathbb{Z}_p$ is a **quadratic residue** (Q.R.) if it has a square root in $\mathbb{Z}_p$

p odd prime $\Rightarrow$ the # of Q.R. in $\mathbb{Z}_p$ is (p-1)/2 + 1

Euler's theorem

Thm: x in $(Z_p)^*$ is a Q.R. $\Longleftrightarrow$ $x^{(p-1)/2} = 1$ in Z_p (p odd prime)

Example:

in $\mathbb{Z}_{11}$:	1^5,	2^5,	3^5,	4^5,	5^5,	6^5,	7^5,	8^5,	9^5,	10^5
=	1	-1	1	1	1,	-1,	-1,	-1,	1,	-1

Note: $x \neq 0 \Rightarrow x^{(p-1)/2} = \left(x^{p-1}\right)^{1/2} = 1^{1/2} \in \{ 1, -1 \}$ in Z_p

Def: $x^{(p-1)/2}$ is called the **Legendre Symbol** of x over p (1798)

Computing square roots mod p

Suppose $p = 3 \pmod 4$

Lemma: if $c \in (Z_p)^*$ is Q.R. then $\sqrt{c} = c^{(p+1)/4}$ in Z_p

Solving quadratic equations mod p

Solve: $a \cdot x^2 + b \cdot x + c = 0 \quad \text{in} \quad Z_p$

Solution: $x = \quad (-b \pm \sqrt{b^2 - 4 \cdot a \cdot c}\,) \;/\; 2a \quad \text{in} \quad Z_p$

- Find $(2a)^{-1}$ in Z_p using extended Euclid.

- Find square root of $b^2 - 4 \cdot a \cdot c$ in Z_p (if one exists)

  using a square root algorithm

Computing e'th roots mod N ??

Let N be a composite number and e>1

When does $c^{1/e}$ **in** Z_N exist? Can we compute it efficiently?

Answering these questions requires the factorization of N

(as far as we know)

Easy problems

- Given composite N and x in Z_N find x^{-1} in Z_N

- Given prime p and polynomial $f(x)$ in $Z_p[x]$

  find x in Z_p s.t. $f(x) = 0$ in Z_p (if one exists)

  Running time is linear in deg(f) .

… but many problems are difficult

Intractable problems with primes

Fix a prime p>2 and g in $(Z_p)^*$ of order q.

Consider the function: $x \mapsto g^x$ in Z_p

Now, consider the inverse function:

$\mathrm{Dlog}_g(g^x) = x$ **where x in {0, …, q-2}**

Example:

in $\mathbb{Z}_{11}$:	1,	2,	3,	4,	5,	6,	7,	8,	9,	10
$\mathrm{Dlog}_2(\cdot)$:	0,	1,	8,	2,	4,	9,	7,	3,	6,	5

Intractable problems with composites

Consider the set of integers: (e.g. for n=1024)

$$\mathbb{Z}_{(2)}(n) \;:=\; \{\, N = p \cdot q \quad \text{where } p,q \text{ are n-bit primes} \,\}$$

<u>Problem 1</u>: Factor a random N in $\mathbb{Z}_{(2)}(n)$ (e.g. for n=1024)

<u>Problem 2</u>: Given a polynomial **f(x)** where degree(f) > 1

and a random N in $\mathbb{Z}_{(2)}(n)$

find x in $\mathbb{Z}_N$ s.t. f(x) = 0 in $\mathbb{Z}_N$

The factoring problem

Gauss (1805): *"The problem of distinguishing prime numbers from composite numbers and of resolving the latter into their prime factors is known to be one of the most important and useful in arithmetic."*

Best known alg. (NFS): run time exp($\tilde{O}(\sqrt[3]{n})$) for n-bit integer

Current world record: **RSA-768** (232 digits)

- Work: two years on hundreds of machines
- Factoring a 1024-bit integer: about 1000 times harder

$\Rightarrow$ likely possible this decade